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## FURTHER LATTICE PACKINGS IN HIGH DIMENSIONS

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*Abstract.* Barnes and Sloane recently described a “general construction” for lattice packings of equal spheres in Euclidean space. In the present paper we simplify and further generalize their construction, and make it suitable for iteration. As a result we obtain lattice packings in  $\mathbb{R}^m$  with density  $\Delta$  satisfying  $\log_2 \Delta \sim -m \log_2^* m$ , as  $m \rightarrow \infty$ , where  $\log_2^* m$  is the smallest value of  $k$  for which the  $k$ -th iterated logarithm of  $m$  is less than 1. These appear to be the densest lattices that have been explicitly constructed in high-dimensional space. New records are also established in a number of lower dimensions, beginning in dimension 96.

§1. *Introduction.* In the past 23 years a series of papers [1–3, 5, 8–12, 15–17] have described a variety of methods for packing equal spheres in Euclidean space. The present paper continues the sequence by simplifying the “general construction” of Barnes and Sloane [1], and thereby eliminating the requirement that the initial lattice be generated by its minimal vectors. This enables us to iterate the construction, obtaining lattice packings in  $\mathbb{R}^m$  with density  $\Delta$  satisfying

$$\log_2 \Delta \sim -m \log_2^* m, \quad \text{as } m \rightarrow \infty, \quad (1)$$

where  $\log_2^* m$  is the smallest value of  $k$  for which the  $k$ -th iterated logarithm of  $m$  is less than 1. These appear to be the densest lattice packings that have been explicitly constructed in high-dimensional space. Non-lattice packings with  $\log_2 \Delta > -6m + o(m)$  were constructed in [15] (and non-lattice packings satisfying (1) in [3]), but there is still room for improvement since the best lattice or non-lattice packings are known to lie in the range

$$-m < \log_2 \Delta < -0.599m + o(m), \quad \text{as } m \rightarrow \infty \quad (2)$$

(see [4, 7, 14, 18]). The new construction also generalizes that given in [1] so as to use codes over other alphabets.

The construction is described in Section 2, and Section 3 contains a number of applications. A selection of the best packings known in dimensions up to  $2^{20}$  will be found in Table 1.

*Notation.* The norm of a vector  $x$  is its squared length  $x \cdot x$ . If  $L_m$  is a lattice in  $\mathbb{R}^m$  (usually the subscript indicates the dimension), its minimum norm  $M$  is  $\min \{x \cdot x : x \in L_m, x \neq 0\}$ , its determinant  $\det L_m$  is the volume of a fundamental region, and its density  $\Delta$  and centre density  $\delta$  are given by

$$\Delta = \frac{V_m(\sqrt{M/2})^m}{\det L_m}, \quad \text{and } \delta = \frac{\Delta}{V_m},$$

where  $V_m$  is the volume of an  $m$ -dimensional sphere of unit radius. In Section 3 we also use the parameters

$$\mu = \frac{1}{2}m, \quad \gamma = \log_2 \delta.$$

All the packings mentioned in this paper are lattice packings.

§2. *The Construction.* The ingredients for the construction are a lattice  $\Lambda$  in  $\mathbb{R}^m$ , an endomorphism  $D$  of  $\Lambda$  that satisfies certain conditions, one of which is that  $\Lambda/D\Lambda$  is an elementary abelian group  $E$  of order  $p^b$ , say, and a family of codes  $C_0 \supseteq C_1 \supseteq \dots \supseteq C_a$  of length  $n$  over  $E$ ; the result is a family  $L_0 \subseteq L_1 \subseteq \dots \subseteq L_a$  of lattice packings in  $\mathbb{R}^{mn}$ .

*Hypotheses.*

- (i) Let  $\Lambda$  be a lattice in  $\mathbb{R}^m$  with minimum non-zero norm  $M$ .
- (ii) Let  $D$  be an endomorphism of  $\Lambda$  which is also a similarity (*i.e.* a constant times an orthogonal transformation) and which satisfies

$$pD^{-1} = \sum_{i=0}^r a_i D^i \tag{3}$$

for integers  $p \geq 1, r \geq 0, a_0, \dots, a_r$ . Let  $T = D^{-1}$ , so that, from (3),  $pT$  is also an endomorphism of  $\Lambda$ .

- (iii) Assume that  $\Lambda/D\Lambda$  has the structure of an elementary abelian group  $E$  of order  $p^b$  for some integer  $b \geq 1$ . This implies that there is a  $b$ -dimensional sublattice  $K \subseteq \Lambda$ , spanned, say, by vectors  $v_1, \dots, v_b \in \Lambda$ , such that

$$K/(D\Lambda \cap K) \cong E.$$

The assumptions also imply that

$$pK \subseteq K \subseteq \Lambda, \tag{4}$$

$$pK \subseteq p\Lambda \subseteq D\Lambda \subseteq \Lambda, \tag{5}$$

and therefore that

$$K/pK = K/(D\Lambda \cap K) \cong E. \tag{6}$$

Note that (4) and (5) imply  $pK \subseteq D\Lambda \cap K$ , so  $K/pK \supseteq K/(D\Lambda \cap K)$ . But

both sides have order  $p^b$  and so must be equal, and (6) follows. Furthermore

$$|\det T| = \frac{1}{p^b} = t^m, \text{ say,} \tag{7}$$

and  $T$  multiplies norms by  $t^2$ .

- (iv) Assume that all  $p^b - 1$  non-zero congruence classes of  $T\Lambda/\Lambda$  have minimum norm at least  $t^2 M$ .
- (v) Let  $\phi$  denote the natural map from  $\mathbb{Z}/p\mathbb{Z}$  to  $\mathbb{Z}$  which takes the congruence class  $\bar{x}$  to  $x$ , for  $x \in \{0, 1, \dots, p-1\}$ . The elements of  $E$  may be identified with the  $b$ -tuples  $\bar{X} = (\bar{x}_1, \dots, \bar{x}_b)$ , where all  $\bar{x}_i \in \mathbb{Z}/p\mathbb{Z}$ . Then

$$\begin{aligned} \bar{X} \rightarrow V(\bar{X}) &= \phi(\bar{x}_1)v_1 + \dots + \phi(\bar{x}_b)v_b \\ &= x_1v_1 + \dots + x_bv_b \end{aligned}$$

is a map from  $E$  into  $\Lambda$ , and

$$\bar{X} = (\bar{X}_1, \dots, \bar{X}_n) \rightarrow V(\bar{X}) = (V(\bar{X}_1), \dots, V(\bar{X}_n))$$

maps  $E^n$  into  $\Lambda^n$ .

- (vi) Let  $C_0 \supseteq C_1 \supseteq \dots \supseteq C_a$  be additive codes† over  $E$  of length  $n$ , where  $C_i$  contains  $p^{bk_i}$  codewords and has minimum distance  $d_i$  (we indicate this by saying that  $C_i$  has parameters  $[n, k_i, d_i]$ ), and suppose that  $C_0$  is the trivial  $[n, n, 1]$  code. Let  $\mathbf{c}_1, \dots, \mathbf{c}_{bk_a} \in E^n$  be chosen so that a typical codeword of  $C_i$  can be written as

$$\sum_{j=1}^{bk_i} \bar{x}_j \mathbf{c}_j, \quad \bar{x}_j \in \mathbb{Z}/p\mathbb{Z},$$

for  $i = 1, \dots, a$ .

*The New Lattices.* Let  $L_0 = \Lambda^n$  and, for  $i = 1, 2, \dots, a$ , define

$$L_i = \bigcup_{x_1, \dots, x_{bk_i}} \left\{ L_{i-1} + \sum_{j=1}^{bk_i} x_j T^i V(\mathbf{c}_j) \right\}, \tag{8}$$

where  $x_1, \dots, x_{bk_i} \in \{0, 1, \dots, p-1\}$ . By abuse of notation we shall also use  $D$  and  $T$  to denote the maps  $(D, D, \dots, D)$  and  $(T, T, \dots, T)$  acting on  $\mathbb{R}^{mn}$ . It is clear that  $L_0$  is a lattice and that  $D$  and  $pT$  are endomorphisms of  $L_0$ .

**THEOREM 1.** For  $i = 1, \dots, a$ ,

- (a)  $L_i$  is a lattice, and in fact

$$L_i = \mathbb{Z}\langle L_{i-1}, T^i V(\mathbf{c}_1), \dots, T^i V(\mathbf{c}_{bk_i}) \rangle; \tag{9}$$

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†I.e. abelian subgroups of  $E^n$  (compare Delsarte [6]).

- (b)  $D$  maps  $L_i$  into  $L_{i-1}$ ; and
- (c)  $pT$  is an endomorphism of  $L_i$ .

*Proof.* The proof is by induction on  $i$ , the results for  $i = 0$  having already been mentioned. (a) To show that  $L_i$  is a lattice, we write any integer combination of the elements on the right-hand side of (8) as

$$l + p \sum_{j=1}^{bk_i} y_j T^i V(\mathbf{c}_j) + \sum_{j=1}^{bk_i} z_j T^i V(\mathbf{c}_j), \tag{10}$$

where  $l \in L_{i-1}$  and  $0 \leq z_j \leq p-1$ . But

$$pT^i V(\mathbf{c}_j) = pT \cdot T^{i-1} V(\mathbf{c}_j) \in L_{i-1}$$

by the induction hypothesis, so (10) becomes

$$l' + \sum_{j=1}^{bk_i} z_j T^i V(\mathbf{c}_j)$$

where  $l' \in L_{i-1}$ , which by (8) is in  $L_i$ . Thus  $L_i$  is a lattice, and therefore can be defined by the right-hand side of (9). Now (b) follows from

$$DT^i V(\mathbf{c}_j) = T^{i-1} V(\mathbf{c}_j) \in L_{i-1}.$$

From (3) and (b) we have

$$pT(L_i) = pD^{-1}(L_i) = \sum_{j=0}^r a_j D^j(L_i) \subseteq L_i,$$

which is (c).

**THEOREM 2.** For  $i = 0, 1, \dots, a$ , the determinant, minimum norm and centre density of  $L_i$  are given by

$$\det L_i = \frac{(\det \Lambda)^n}{p^{b \sum_{j=1}^i k_j}}, \tag{11}$$

$$\bar{M} = \min \{M, d_j t^{2j} M \text{ for } j = 1, \dots, i\}, \tag{12}$$

and

$$\delta = \bar{M}^{mn/2} / 2^{mn} \det L_i \tag{13}$$

respectively.

*Proof.* Equation (11) is immediate from (8), and (12) follows from the definition of  $L_i$  and the fact that the minimum distance of  $C_i$  is  $d_i$ . For more details compare the proof of Theorem 3 in [1].

Usually we are only interested in the finest lattice  $L_a$ . The construction may now be applied to  $L_a$ , since it inherits  $D$  and  $T$  from  $\Lambda$ , and (3) still holds. The values of  $p$  and  $t$  are unchanged, while  $b$  becomes  $nb$ .

§3. *Examples.* The above construction includes some of the most important special cases of Constructions *A* and *B* of [12] (see [1]), as well as the constructions of [1], and so a large number of examples are already known. In this section we give several additional examples of some interest.

If *D* is a norm-doubler. In all these examples *D* will be a norm-doubler, so that  $t = 1/\sqrt{2}$ ,  $p = 2$  and  $b = m/2$ . We shall take the codes  $C_i, i \geq 1$ , to be maximal distance separable codes over the field  $GF(2^b)$ , with parameters  $[n, k_i = n - 2^i + 1, d_i = 2^i]$ . In general, the largest  $n$  for which such codes are presently known to exist is  $2^b + 1$  (see [13, Chapter 11]). When  $n = 2^b + 1, C_i, i \geq 1$ , can be taken to be the cyclic code with generator polynomial

$$g(x) = \prod_{j=-s}^s (x + \xi^j),$$

where  $\xi \in GF(2^{2b})$  is a primitive  $n$ -th root of unity and  $s = 2^{i-1} - 1$  (see the proof of Theorem 9, Chapter 11 of [13])†. Besides being additive, these codes are also closed under multiplication by elements of  $GF(2^b)$ . However we make no use of this multiplicative structure in our construction. Clearly  $C_0 \supseteq C_1 \supseteq C_2 \supseteq \dots$ . For  $n$  less than  $2^b + 1$  the codes are shortened by setting the appropriate number of information symbols equal to zero. If  $n$  lies in the range  $2 \leq n \leq 2^{b+1}$ , we use the codes  $C_1, C_2, \dots, C_a$ , where  $a$  is determined by  $2^a \leq n < 2^{a+1}$ . Then the value of  $\sum_{j=1}^a k_j$  for use in (11) is

$$an - 2^{a+1} + a + 2. \tag{14}$$

Since  $d_j t^{2j} = 1$ , the minimum norm is unchanged (see (12)). This version of the construction therefore converts a lattice  $\Lambda$  in  $\mathbb{R}^m$  with centre density  $\delta = 2^v$  (say) into a lattice  $\Lambda' = L_a$  in  $\mathbb{R}^{mn}$  with centre density  $\delta' = 2^v$ , where from (11), (13), (14)

$$\gamma' = n\gamma + \frac{m}{2}(an - 2^{a+1} + a + 2). \tag{15}$$

Our examples are all descendants of the three lattices  $\mathbb{Z}^2, \Lambda_{24}$  and  $\mathbb{P}48q$ .

*Packings constructed from the Lattice  $\mathbb{Z}^2$ .* For the first example we take  $\Lambda$  to be the familiar two-dimensional square lattice  $\mathbb{Z}^2$  with minimum norm  $M = 1$  (see Figure 1). Although this is not a particularly dense packing, it has surprisingly good progeny. We let *D* map (1, 0) to (1, 1), and (0, 1) to (1, -1), or in matrix notation, with *D* mapping  $v$  to  $vD$ ,

$$D = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, \quad T = D^{-1} = \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix},$$

so that (Eq. 3) reads  $2D^{-1} = D$ . Then  $\Lambda/D\Lambda \cong \mathbb{Z}/2\mathbb{Z}, p^b = 2^1$ , and we may take *K* to be the 1-dimensional lattice spanned by  $v_1 = (1, 0)$  (see Figure 1).

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† That theorem is incorrect if the field size  $q$  is odd (see [13, p. xii]) but that does not concern us since here  $q$  is even.

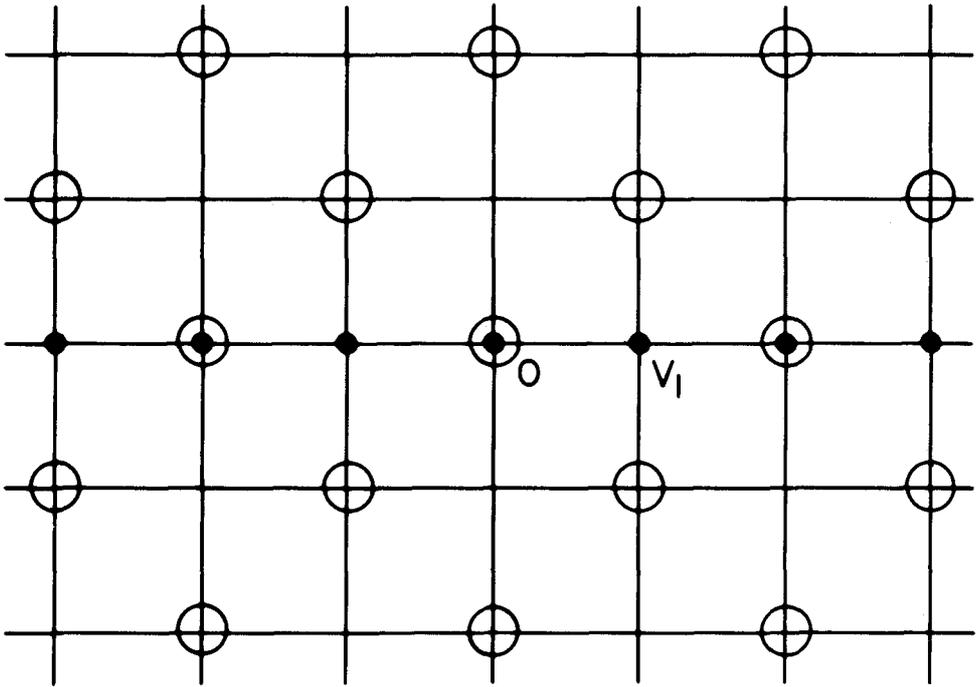


Figure 1. The lattice  $\Lambda = \mathbb{Z}^2$  (represented by the underlying grid of points), and the sublattices  $D\Lambda$  (small circles) and  $K$  (solid circles).

We apply our construction to  $\mathbb{Z}^2$ , with  $n = 2$  and  $C_1$  equal to the code  $\{00, 11\}$  over  $GF(2)$ , and obtain the Schläfli lattice  $D_4$ , with  $\gamma = -3$  (from (15)).

From now on we need not specify  $D$ ,  $T$  or  $p^b$ , since these are determined (see the remarks at the end of Section 2). Applying the construction to  $D_4$  with  $n = 1, \dots, 5$  we obtain the laminated lattices (see [5])

$$\Lambda_4 = D_4, \quad \Lambda_8 = E_8, \quad \Lambda_{12}, \quad \Lambda_{16}, \quad \Lambda_{20}$$

for which  $\gamma$  is

$$-3, -4, -5, -4, -3$$

respectively.

$E_8$  (the first grandchild of  $\mathbb{Z}^2$ ) was one of the examples studied in [1]†. It produces  $\Lambda_{16}$  again, packings  $\bar{\Lambda}_{32}$  and  $\bar{\Lambda}_{40}$  with the same density as the laminated lattice  $\Lambda_{32}$  and  $\Lambda_{40}$ , excellent packings in dimensions 80 to 136, but indifferent packings in dimensions 24 and 48 to 72.

Table 1 contains a selection of the best lattice packings known up to dimension  $2^{20}$ . The second column gives  $\gamma = \log_2 \delta$ , while a parenthesized entry ( $\Lambda$ ) in the third column indicates that the packing can be obtained by applying our construction to the lattice  $\Lambda$ . Those packings (in dimension  $\geq 96$ ) for which no reference is given are believed to set new records for lattices.

The children of  $E_8$  extend to dimension 136 in the table and then stop. Beyond this point the lattices obtained by applying the construction to  $\Lambda_{12}$  and  $\Lambda_{16}$  are roughly comparable, up to dimension  $16(2^8 + 1)$ . The children of  $\Lambda_{20}$ ,  $\bar{\Lambda}_{32}$  and  $\bar{\Lambda}_{40}$

† The map  $D$  used in [1] differs from the one given here, but is equivalent to it. One can in fact show that there is essentially only one norm-doubling map for  $E_8$  with the requisite properties.

Table 1. A selection of the densest known packings  
 $n$  = dimension,  $\gamma$  =  $\log_2$  (centre density)

$n$	$\gamma$	name	Ref.	$n$	$\gamma$	name	Ref.
0	0	$\Lambda_0$		492	683	$(\Lambda_{12})$	
4	-3	$D_4$		496	660	$(\Lambda_{16})$	
8	-4	$E_8$		500	665	$(\Lambda_{20})$	
16	-4	$\Lambda_{16}$	[2]	504	708	$(\Lambda_{12})$	
24	0	$\Lambda_{24}$	[9]	512	698	$B_{512}$	[1]
32	0	$\Lambda_{32}$	[2,12]	516	733	$(\Lambda_{12})$	
40	4	$\Lambda_{40}$	[12]	1024	1856	$(\Lambda_{16})$	
48	14.039	$\mathbb{P}48q$	[12]	1032	1896	$(\Lambda_{24})$	[1]
64	22	—	[1]	1040	1900	$(\Lambda_{16})$	
80	36	$(E_8)$	[1]	2040	4680	$(\Lambda_{24})$	[1]
96	52.078	$(\mathbb{P}48q)$		2048	4680	$(\Lambda_{16})$	
128	88	$(E_8)$	[1]	2064	4752	$(\Lambda_{24})$	[1]
136	100	$(E_8)$	[1]	4096	11344	$(\Lambda_{16})$	
160	112	$(\Lambda_{16})$		4104	11400	$(\Lambda_{24})$	[1]
168	120	$(\Lambda_{24})$	[1]	8208	26808	$(\Lambda_{24})$	[1]
176	132	$(\Lambda_{16})$		16392	61608	$(\Lambda_{24})$	[1]
192	156	$(\Lambda_{24})$	[1]	32784	139488	$(\Lambda_{24})$	[1]
256	250	$B_{256}$	[1]	65544	311496	$(\Lambda_{24})$	[1]
264	264	$(\Lambda_{24})$	[1]	131088	$6.613 \times 10^5$	$(\mathbb{P}48q)$	
272	268	$(\Lambda_{16})$		262176	$1.453 \times 10^6$	$(\mathbb{P}48q)$	
288	300	$(\Lambda_{24})$	[1]	524304	$3.168 \times 10^6$	$(\mathbb{P}48q)$	
384	464	$(\Lambda_{16})$		1048608	$6.861 \times 10^6$	$(\mathbb{P}48q)$	

are not quite as good as other known packings. However, examination of Table 1 (and of a more extensive table of which this is an extract) shows that at the moment one cannot completely ignore the children of any of these lattices.

*Packings constructed from the Leech lattice  $\Lambda_{24}$ .* The Leech lattice  $\Lambda_{24}$  is another example that was studied in [1], using the norm-doubling map  $D = I - i$ , where  $i$  is a certain automorphism of  $\Lambda_{24}$ . There is a second choice for  $D$ , essentially different from  $I - i$ , which may be defined as follows. Let the 24 coordinate positions be labelled  $\infty, 0, 1, \dots, 22$ , and as usual (see [9, 12]) let  $\Lambda_{24}$  be spanned by the vectors:—

- $(2^{12}, 0^{12})$ , 23 vectors, supported on a translate of  $\{0\} \cup q$ , where  $q$  denotes the set of non-zero quadratic residues modulo 23,
- $(-3, 1^{23})$ , a single vector, and
- $(\pm 4^2, 0^{22})$ ,  $4 \cdot \binom{24}{2}$  vectors.

Then we may take  $D$  to be represented by the matrix

$$\frac{1}{4} \begin{vmatrix} 3 & -1 & -1 & -1 & -1 & -1 & \dots & -1 \\ \hline -1 & -3 & 1 & 1 & 1 & 1 & \dots & -1 \\ -1 & 1 & 1 & 1 & 1 & -1 & \dots & -3 \\ -1 & 1 & 1 & 1 & -1 & 1 & \dots & 1 \\ & & \dots & & & & \dots & \\ -1 & -1 & -3 & 1 & 1 & 1 & \dots & -1 \end{vmatrix}$$

where in the second row there is a  $-3$  in position  $0$ ,  $1$ 's in the positions in  $q$ , and  $-1$ 's elsewhere. Also  $T = D^{-1} = \frac{1}{2}D$  and  $p^b = 2^{12}$ . Applying our construction to  $\Lambda_{24}$  using either version of  $D$  produces lattices which run neck and neck with  $\Lambda_{12}$  and  $\Lambda_{16}$  up to dimension  $4096$ , and beyond this point are the best packings known up to dimension  $24(2^{12} + 1)$  (see Table 1 and [1]).

*Packings constructed from  $\mathbb{P}48q$ .* Our third starting point is the 48-dimensional lattice  $\mathbb{P}48q$  constructed in [12], which has a norm-doubling map very similar to the second map found for  $\Lambda_{24}$ . If the coordinate positions are labelled  $\infty, 0, 1, \dots, 46$ , we may take  $\mathbb{P}48q$  to be spanned by the vectors:—

- $(2^{24}, 0^{24})$ , 47 vectors, supported on a translate of  $\{0\} \cup q$ , where  $q$  denotes the set of non-zero quadratic residues modulo 47,
- $(-5, 1^{47})$ , a single vector, and
- $(\pm 6^2, 0^{46})$ ,  $4 \cdot \binom{48}{2}$  vectors.

Then we may take  $D$  to be

$$\frac{1}{6} \begin{vmatrix} 5 & -1 & -1 & -1 & -1 & -1 & \dots & -1 \\ \hline -1 & -5 & 1 & 1 & 1 & 1 & \dots & -1 \\ -1 & 1 & 1 & 1 & 1 & -1 & \dots & -5 \\ -1 & 1 & 1 & 1 & -1 & 1 & \dots & 1 \\ & & & & \dots & & & \dots \\ -1 & -1 & -5 & 1 & 1 & 1 & \dots & -1 \end{vmatrix}$$

where in the second row there is a  $-5$  in position  $0$ ,  $1$ 's in the positions in  $q$ , and  $-1$ 's elsewhere. Also  $T = D^{-1} = \frac{1}{2}D$ . This gives excellent packings in dimensions 96 and from  $24(2^{12} + 1)$  to  $48(2^{24} + 1)$ .

*Higher dimensions.* The construction may be applied repeatedly to any of these packings, producing an infinite tree of lattices. We may derive a lower bound to the density obtained in this way in high dimensions by the following argument. Let  $\Lambda$  be any of the lattices described in this section, with say dimension  $m = 2\mu$ , centre density  $\delta = 2^\gamma$ , and norm-doubling map  $D$ . We apply our construction using codes of length  $n = 2^\mu$ , obtaining a lattice  $\Lambda' (= L_a)$  in dimension  $m' = 2\mu'$  with centre density  $\delta' = 2^{\gamma'}$ , where, from (15)

$$\mu' = \mu 2^\mu, \tag{16}$$

$$\gamma' = 2^\mu \gamma + \mu(\mu - 2)2^\mu + \mu(\mu + 2). \tag{17}$$

It is simpler to work with the density  $\Delta$  rather than the centre density  $\delta$ , so let us define

$$\eta = -\frac{\log_2 \Delta}{m}.$$

(We know from (2) that for large dimensions  $m$  the best packings satisfy  $0.599 \leq \eta \leq 1$ .) Now

$$\log_2 \Delta = \log_2 \delta - \mu \log_2 \mu + \mu \log_2 (\pi e) - \frac{1}{2} \log_2 (2\pi\mu) + o(1),$$

and so (17) becomes

$$\eta' = \eta + 1 - \frac{1}{4\mu} \log_2 (2\pi\mu) + o\left(\frac{1}{\mu}\right). \tag{18}$$

The solution of (16) and (18) is

$$\eta(m) = \log_2^*(m)$$

and therefore applying the construction repeatedly leads to lattices satisfying (1). Note that this asymptotic behavior is independent of the choice of the initial lattice.

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