

## Codes over $GF(4)$ and Complex Lattices

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This paper studies the relationship between error-correcting codes over  $GF(4)$  and complex lattices (more precisely,  $\mathbb{Z}[\omega]$ -modules in  $\mathbb{C}^n$ , where  $\omega = e^{2\pi i/3}$ ). The theta-functions of self-dual lattices are characterized. Two general methods are presented for constructing lattices from codes. Several examples are given, including a new lattice sphere-packing in  $\mathbb{R}^{36}$ .

### 1. INTRODUCTION

The connections between binary and ternary error-correcting codes on the one hand, and lattices and sphere-packings in  $\mathbb{R}^n$  on the other have been studied by several authors [6, 7, 19-21, 26, 29, 30, 39, 40]. Since the theorem of Gleason, Pierce and Turyn [1, p. 381; 2] states that codes over  $GF(2)$ ,  $GF(3)$  and  $GF(4)$  have especially nice properties not shared by codes over other fields, one might expect that codes over  $GF(4)$  should also be connected with sphere-packings. It is the aim of the present paper to show that there are natural connections between codes over  $GF(4)$  and complex lattices and sphere-packings. For example the weight enumerator of the code and the theta-function of the lattice have similar properties: compare Theorem 3 below with Theorem 1, and Theorem 4 with Theorem 9. The cusp form  $\Delta(z) = q^2 \prod_{m=1}^{\infty} (1 - q^{2m})^{24} = \sum_{m=0}^{\infty} a_m q^{2m}$  (where the  $a_m$  are the Ramanujan numbers and  $q = e^{\pi iz}$ ) plays a special role in characterizing the theta-functions of even unimodular lattices in  $\mathbb{R}^n$ . The same role for our complex lattices is played by  $\Delta_6(z)$ ; compare Eqs. (19), (20), and Theorems 9, 10. In Sections 6 and 8 two constructions are given for obtaining lattices from codes over  $GF(4)$ . In this way it is possible to construct the densest known lattice packings  $A_2$  in  $\mathbb{R}^2$ ,  $E_6$  in  $\mathbb{R}^6$  and  $K_{12}$  in  $\mathbb{R}^{12}$  directly from codes (filling in a gap in [21]); we also find a very dense lattice packing in  $\mathbb{R}^{36}$  (Example 7).

Complex lattices have also been used in studying complex polytopes and simple groups [8-11, 14, 14a, 23, 36, 37] and are related to the problem of finding sets of lines in  $\mathbb{C}^n$  having a prescribed number of angles [12]

2. COMPLEX LATTICES AND SPHERE-PACKINGS

A (real) lattice in  $\mathbb{R}^n$  is a discrete subgroup of  $\mathbb{R}^n$  which spans  $\mathbb{R}^n$ , in the sense that any point of  $\mathbb{R}^n$  can be written as a linear combination of lattice points with real coefficients. In particular a real lattice is a  $\mathbb{Z}$ -module. In this paper a *complex lattice* (or  $\mathcal{E}$ -lattice) in  $\mathbb{C}^n$  will always mean a discrete subgroup of  $\mathbb{C}^n$  which (i) spans  $\mathbb{C}^n$  in the sense that any point of  $\mathbb{C}^n$  can be written as a linear combination of lattice points with complex coefficients, and (ii) is a free  $\mathcal{E}$ -module, where  $\mathcal{E} = \{a + b\omega : a, b \in \mathbb{Z}\}$ ,  $\omega = e^{2\pi i/3}$ , denotes the Eisenstein integers (cf. [9, p. 421; 10, p. 145; 11; 14, p. 635; 17, p. 179; 22]). Thus  $\mathcal{E}$  plays the same role for these complex lattices as  $\mathbb{Z}$  plays for real lattices. In particular, the dual lattice  $\Lambda^\perp = \{\mathbf{x} \in \mathbb{C}^n : \mathbf{x} \cdot \bar{\mathbf{y}} \in \mathcal{E} \text{ for all } \mathbf{y} \in \Lambda\}$ , where  $\mathbf{x} \cdot \mathbf{y} = \sum_{i=1}^n x_i y_i$  and the bar denotes conjugation.

Let  $\{\mathbf{e}^1 = (e_1^1, \dots, e_n^1), \dots, \mathbf{e}^n = (e_1^n, \dots, e_n^n)\}$  be an  $\mathcal{E}$ -basis for  $\Lambda$ . The  $n \times n$  matrix  $M$  with  $(i, j)$ th entry equal to  $e_j^i$  is called a *generator matrix* for  $\Lambda$ . Any  $\mathbf{x} \in \Lambda$  can be written uniquely as  $\mathbf{x} = \mathbf{c}M$  where  $\mathbf{c} \in \mathcal{E}^n$ . It is easy to see that  $(M^{-1})^{tr}$  is a generator matrix for  $\Lambda^\perp$ . The determinant of  $\Lambda$  is defined to be

$$\det \Lambda = |\det M|,$$

and  $\det \Lambda^\perp = (\det \Lambda)^{-1}$ . If  $\Lambda = \Lambda^\perp$  then  $\Lambda$  is called *self-dual*. For example  $\mathcal{E}^n$  is a self-dual lattice in  $\mathbb{C}^n$  with  $M = I_n$ . Also  $\Lambda$  is self-dual if and only if  $\det \Lambda = 1$  and  $\mathbf{e}^i \cdot \bar{\mathbf{e}}^j \in \mathcal{E}$  for all  $i, j$ . The self-dual lattices in  $\mathbb{C}^n$  for  $n \leq 12$  have been classified by Feit [14a].

The *minimum squared length* of  $\Lambda$  is

$$d = \min_{\mathbf{x} \in \Lambda - \{0\}} \mathbf{x} \cdot \bar{\mathbf{x}}.$$

If complex spheres of radius  $\rho = \frac{1}{2}d^{1/2}$  are centered at the points of  $\Lambda$  these spheres do not overlap and form a complex  $\mathcal{E}$ -lattice sphere-packing in  $\mathbb{C}^n$ . The *density* of this packing or lattice is the fraction of  $\mathbb{C}^n$  covered by the spheres, given by

$$\Delta = \frac{V_{2n} \rho^{2n}}{(\det \Lambda)^2 (3^{1/2}/2)^n}, \tag{1}$$

where  $V_n = \pi^{n/2} / \Gamma(n/2 + 1)$ . The *contact number*, or number of spheres touching one sphere, is  $\tau = \#\{\mathbf{x} \in \Lambda : \mathbf{x} \cdot \bar{\mathbf{x}} = d\}$ . Although it is not our goal in the present paper, one can now study  $\Delta^o(n)$  and  $\tau^o(n)$ , the largest  $\Delta$  and  $\tau$  attained by  $\mathcal{E}$ -lattice packings in  $\mathbb{C}^n$ . The determinant, density and contact number of a real lattice are defined in the analogous way (see [34, 21]).

If  $\Lambda$  is an  $\mathcal{E}$ -lattice in  $\mathbb{C}^n$  with parameters  $d, \Delta, \tau$ , let  $\Lambda_R \subset \mathbb{R}^{2n}$  consist of all points

$$(\operatorname{Re}(x_1), \operatorname{Im}(x_1), \operatorname{Re}(x_2), \dots, \operatorname{Im}(x_n))$$

where  $(x_1, \dots, x_n) \in \Lambda$ . Then  $\Lambda_R$  is a lattice in  $\mathbb{R}^{2n}$  with minimum squared length  $d$ . If real spheres of radius  $\rho = \frac{1}{2}d^{1/2}$  are centered at the points of  $\Lambda_R$ , a lattice sphere-packing in  $\mathbb{R}^{2n}$  is obtained, having determinant  $\det \Lambda_R = (3^{1/2}/2)^n (\det \Lambda)^2$ , density  $= \Delta$ , and contact number  $= \tau$ . (Note that in general  $(\Lambda^\perp)_R \neq (\Lambda_R)^\perp$ .) Since the largest  $\Lambda$  for a lattice packing in  $\mathbb{R}^n$  is known for  $n \leq 8$  (see [34, 21]), we have

$$\begin{aligned} \Delta^o(1) &\leq V_2 \cdot 2^{-1} \cdot 3^{-1/2}, & \Delta^o(2) &\leq V_4 \cdot 2^{-3}, & \Delta^o(3) &\leq V_6 \cdot 2^{-3} \cdot 3^{-1/2}, \\ & & \Delta^o(4) &\leq V_8 \cdot 2^{-4} & & \end{aligned} \quad (2)$$

Examples 4 and 5 will show that equality holds for  $\Delta^o(1)$  and  $\Delta^o(3)$ .

### 3. THE THETA-FUNCTION OF A LATTICE

The theta-function of an  $\mathcal{E}$ -lattice  $\Lambda$  in  $\mathbb{C}^n$  is

$$\Theta_\Lambda(z) = \sum_{\mathbf{x} \in \Lambda} e^{\pi i z \mathbf{x} \cdot \bar{\mathbf{x}}}$$

and is holomorphic for  $z \in \mathcal{H} = \{z \in \mathbb{C} : \text{Im}(z) > 0\}$ .

**THEOREM 1.** For  $z \in \mathcal{H}$ ,

$$\Theta_{\Lambda^\perp}(z) = (\det \Lambda)^2 (2i/3^{1/2}z)^n \Theta_\Lambda(-4/3z). \quad (3)$$

*Remark.* For a real lattice  $\Lambda \subset \mathbb{R}^n$  with theta-function  $\Theta_\Lambda(z) = \sum_{\mathbf{x} \in \Lambda} \times \exp(\pi i z \mathbf{x} \cdot \mathbf{x})$ , the analogous result is

$$\Theta_{\Lambda^\perp}(z) = (\det \Lambda)(i/z)^{n/2} \Theta_\Lambda(-1/z). \quad (4)$$

(see for example [13, Chap. 2, Sect. 11.3; 35, Chap. VII, Sect. 6].) Note that an  $\mathcal{E}$ -lattice  $\Lambda$  and the real lattice  $\Lambda_R$  have the same theta-function.

*Proof.* Equation (3), like (4), is a consequence of the appropriate version of the Poisson summation formula. For  $z = z_a + \omega z_b = z_r + iz_i \in \mathbb{C}$ , where  $z_a, z_b, z_r, z_i$  are real, define  $R: \mathbb{C} \rightarrow \mathbb{R}$  by

$$R(z) = z_a = z_r + (1/3^{1/2}) z_i.$$

If  $z' = z'_a + \omega z'_b$  then  $R(z\bar{z}')$   $= z_a z'_a - z_a z'_b + z_b z'_b$ . If  $\Lambda$  is an  $\mathcal{E}$ -lattice in  $\mathbb{C}^n$ , the functions  $\chi_{\mathbf{y}}: \mathbb{C}^n/\Lambda \rightarrow \mathbb{C}$  given by

$$\chi_{\mathbf{y}}(\mathbf{x}) = e^{2\pi i R(\mathbf{x} \cdot \bar{\mathbf{y}})}, \quad \mathbf{x} \in \mathbb{C}^n/\Lambda,$$

for  $\mathbf{y} \in \Lambda^\perp$  form the character group of  $\mathbb{C}^n/\Lambda$ , which is isomorphic to  $\Lambda^\perp$ . The

*Poisson summation formula* [13, p. 220; 18, p. 44; 24, p. 153] implies that if  $f$  is a Schwartz function on  $\mathbb{C}^n$  then

$$\sum_{\mathbf{x} \in \Lambda} f(\mathbf{x}) = (\det \Lambda)^{-2} \left(\frac{2}{3^{1/2}}\right)^n \sum_{\mathbf{y} \in \Lambda^\perp} f(\mathbf{y}), \tag{5}$$

where

$$f(\mathbf{y}) = \int_{\mathbb{C}^n} f(\mathbf{x}) e^{-2\pi i R(\mathbf{x} \cdot \bar{\mathbf{y}})} d\mathbf{x}.$$

Equation (3) follows from (5) by taking

$$f(\mathbf{x}) = (4i/3z)^n e^{-4\pi i \mathbf{x} \cdot \mathbf{x}/3z}, \quad f(\mathbf{y}) = e^{\pi i z \bar{\mathbf{y}} \cdot \bar{\mathbf{y}}}.$$

**COROLLARY 2.** *If  $\Lambda$  is a self-dual lattice in  $\mathbb{C}^n$ ,*

- (i)  $\theta_\Lambda(z + 2) = \theta_\Lambda(z)$ ,
- (ii)  $\theta_\Lambda(-4/3z) = (3^{1/2}z/2i)^n \theta_\Lambda(z)$ , (6)
- (iii) *if  $n$  is even,  $\theta_\Lambda(2z)$  is a modular form of weight  $n$  for the congruence group  $\Gamma_0(3)$ . (Here we are using "weight" as defined in [35] or [38], not as in [15].)*

*Proof.* (ii) follows from Theorem 1. Put  $f(z) = \theta_\Lambda(2z)$ . Then  $f(z + 1) = f(z)$  and  $f(-1/3z) = (-i 3^{1/2}z)^n f(z)$ , which imply  $f(z/(3z + 1)) = (3z + 1)^n f(z)$ . Since  $\begin{pmatrix} 11 \\ 01 \end{pmatrix}$  and  $\begin{pmatrix} 10 \\ 31 \end{pmatrix}$  generate  $\Gamma_0(3)$ , (iii) follows. Note that (iii) is strictly weaker than (i) and (ii).

#### 4. JACOBI THETA-FUNCTIONS

The theta-functions of lattices are conveniently expressed in terms of the Jacobi theta-functions, defined as follows (cf. [4, 33, 41, 43]).

$$\begin{aligned} \theta'_1(z) &= \theta'_1(0 | z) = 2 \sum_{m=0}^{\infty} (-1)^m (2m + 1) q^{(m+1/2)^2} \\ &= 2q^{1/4} \prod_{m=1}^{\infty} (1 - q^{2m})^3, \end{aligned} \tag{7}$$

$$\theta_2(z) = \theta_2(0 | z) = 2 \sum_{m=0}^{\infty} q^{(m+1/2)^2},$$

$$\theta_3(z) = \theta_3(0 | z) = 1 + 2 \sum_{m=1}^{\infty} q^{m^2},$$

$$\theta_4(z) = \theta_4(0 | z) = 1 + 2 \sum_{m=1}^{\infty} (-q)^{m^2},$$

where  $q = e^{\pi iz}$ ,  $z \in \mathcal{H}$ . The following are the most useful of the "labyrinth" of identities relating the  $\theta_i$ .

$$\theta'_1(z) = \theta_2(z) \theta_3(z) \theta_4(z), \quad (8)$$

$$\theta_2\left(-\frac{1}{z}\right) = \left(\frac{z}{i}\right)^{1/2} \theta_4(z), \quad \theta_3\left(-\frac{1}{z}\right) = \left(\frac{z}{i}\right)^{1/2} \theta_3(z), \quad (9)$$

$$\theta_2(z)^2 = 2\theta_2(2z) \theta_3(2z), \quad \theta_2(z) \theta_4(z) = \theta_4(2z)^2, \quad (10)$$

$$\theta_3(z)^2 + \theta_4(z)^2 = 2\theta_3(2z)^2, \quad \theta_3(z)^2 - \theta_4(z)^2 = 2\theta_2(2z)^2, \quad (11)$$

$$\theta_3(z) + \theta_4(z) = 2\theta_3(4z), \quad \theta_3(z) - \theta_4(z) = 2\theta_2(4z). \quad (12)$$

EXAMPLE 1. The lattice  $\mathcal{E}$  in  $\mathbb{C}^1$  has theta-function

$$\begin{aligned} \Theta_{\mathcal{E}}(z) &= \sum_{a+b\omega \in \mathcal{E}} q^{|a+b\omega|} = \sum_{a=-\infty}^{\infty} \sum_{b=-\infty}^{\infty} q^{a^2-ab+b^2} \\ &= \theta_2(z) \theta_2(3z) + \theta_3(z) \theta_3(3z) = \phi_0\left(\frac{1}{2}z\right) \quad (\text{say}) \\ &= 1 + 6q + 6q^3 + 6q^4 + 12q^7 + 6q^9 + \dots \end{aligned} \quad (13)$$

### 5. CODES OVER $GF(4)$

Let the elements of  $GF(4)$  be  $0, 1, \omega, \omega^2$ , with  $\omega^2 + \omega + 1 = 0$ . The conjugate of  $x \in GF(4)$  is  $\bar{x} = x^2$ . Let  $\mathcal{C}$  be a linear code over  $GF(4)$  of length  $n$  and dimension  $k$  (see [27], [28]). The dual code  $\mathcal{C}^\perp = \{\mathbf{u} \in GF(4)^n: \mathbf{u} \cdot \bar{\mathbf{v}} = 0 \text{ for all } \mathbf{v} \in \mathcal{C}\}$ , where  $\mathbf{u} \cdot \bar{\mathbf{v}} = \sum_{i=1}^n u_i \bar{v}_i$ , and has dimension  $n - k$ . (A different definition of  $\mathcal{C}^\perp$  was used in [27].) If  $\mathcal{C} = \mathcal{C}^\perp$ ,  $\mathcal{C}$  is called *self-dual* (and  $n$  is even). A code can be defined by giving its *generator matrix*  $M$ , which is a  $k \times n$  matrix whose rows span the code. If  $M = [I | B]$  is a generator matrix for  $\mathcal{C}$  then  $[\bar{B}^{tr} | I]$  is a generator matrix for  $\mathcal{C}^\perp$ . Also  $\mathcal{C}$  is self-dual if and only if  $B\bar{B}^{tr} = I$  (i.e.  $B$  is unitary). The weight  $wt(\mathbf{x})$  of a vector  $\mathbf{x} = (x_1, \dots, x_n)$  is the number of nonzero  $x_i$ . The *minimum distance* of  $\mathcal{C}$  is

$$\text{dist}(\mathcal{C}) = \min\{wt(\mathbf{u}): \mathbf{u} \in \mathcal{C}, \mathbf{u} \neq \mathbf{0}\}$$

and the *weight enumerator* is

$$\begin{aligned} W_{\mathcal{C}}(x, y) &= \sum_{\mathbf{u} \in \mathcal{C}} x^{n-wt(\mathbf{u})} y^{wt(\mathbf{u})} \\ &= \sum_{r=0}^n A_r x^{n-r} y^r, \end{aligned}$$

where  $A_r$  is the number of codewords in  $\mathcal{C}$  of weight  $r$ . The automorphism group of  $\mathcal{C}$ ,  $\text{Aut}(\mathcal{C})$ , consists of all  $n \times n$  monomial matrices over  $GF(4)$  which set-wise preserve the code.

THEOREM 3 (MacWilliams [25, 28]).

$$W_{\mathcal{C}^\perp}(x, y) = 4^{-k}W_{\mathcal{C}}(x + 3y, x - y)$$

THEOREM 4 [26]. If  $\mathcal{C}$  is self-dual then  $W_{\mathcal{C}}(x, y) \in \mathbb{C}[f_2, f_6]$ , where

$$f_2 = x^2 + 3y^2, \quad f_6 = y^2(x^2 - y^2)^2.$$

EXAMPLE 2. The self-dual code  $C_2 = \{00, 11, \omega\omega, \omega^2\omega^2\}$  has  $n = 2$ ,  $k = 1$ ,  $\text{dist}(C_2) = 2$ ,  $W_{C_2}(x, y) = f_2$ , and  $\text{Aut}(C_2) = 3\mathcal{S}_2$ .

EXAMPLE 3. The matrix

$$M = \begin{bmatrix} 1 & 0 & 0 & 1 & \omega & \omega \\ 0 & 1 & 0 & \omega & 1 & \omega \\ 0 & 0 & 1 & \omega & \omega & 1 \end{bmatrix}$$

generates a self-dual code  $Q_6$  with  $n = 6$ ,  $k = 3$ ,  $W_{Q_6}(x, y) = x^6 + 45x^2y^4 + 18y^6 = f_2^3 - 9f_6$ , and  $\text{Aut}(Q_6) = 3\mathcal{A}_6$ .

For further properties and examples see [27].

## 6. CONSTRUCTION OF LATTICES FROM CODES OVER $GF(4)$

The starting point is the observation that  $2\mathcal{E}$  is a prime ideal in  $\mathcal{E}$ , and  $\mathcal{E}/2\mathcal{E}$  is isomorphic to  $GF(4)$ . In fact, if we take coset representatives  $0, 1, \omega, \omega^2$  for  $2\mathcal{E}$  in  $\mathcal{E}$ , there is an isomorphism mapping  $2\mathcal{E}$  onto  $0$ ,  $1 + 2\mathcal{E}$  onto  $1$ ,  $\omega + 2\mathcal{E}$  onto  $\omega$ , and  $\omega^2 + 2\mathcal{E}$  onto  $\omega^2$ . Hence there is a map  $\sigma: \mathcal{E} \rightarrow \mathcal{E}/2\mathcal{E} \rightarrow GF(4)$  which sends  $\omega^r \in \mathcal{E}$  onto  $\omega^r \in GF(4)$  ( $r = 0, 1, 2$ ).

Two constructions will be given (cf. [21]), the second in Section 8.

CONSTRUCTION A. Let  $\mathcal{C}$  be a linear code over  $GF(4)$  of length  $n$  and dimension  $k$ . Let  $\Lambda(\mathcal{C})$  consist of all  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{C}^n$  such that  $\sigma(2^{1/2}\mathbf{x}) \in \mathcal{C}$ .

In words, the points of  $\Lambda(\mathcal{C})$  are all vectors  $\mathbf{x}$  that can be obtained from codewords by adding twice an Eisenstein integer to each component, and then dividing by  $2^{1/2}$ .

THEOREM 5. (i)  $\Lambda(\mathcal{C})$  is an  $\mathcal{E}$ -lattice in  $\mathbb{C}^n$ . (ii)  $\det \Lambda = 2^{n/2-k}$ . (iii)  $\Lambda(\mathcal{C}^\perp) = \Lambda(\mathcal{C})^\perp$ . (iv)  $\Lambda(\mathcal{C})$  is self-dual if and only if  $\mathcal{C}$  is a self-dual code.

*Proof.* (i) is immediate. Without loss of generality let  $\mathcal{C}$  have generator matrix  $[I_k | B]$ . Then

$$M = \frac{1}{2^{1/2}} \begin{bmatrix} I_k & B \\ 0 & 2I_{n-k} \end{bmatrix}$$

is a generator matrix for  $\Lambda(\mathcal{C})$ . The rest of the proof is now straightforward.

**THEOREM 6.** *The lattice  $\Lambda(\mathcal{C})$  has the following properties:*

$$d = \begin{cases} \frac{1}{2} \text{dist}(\mathcal{C}) & \text{if } \text{dist}(\mathcal{C}) \leq 4, \\ 2 & \text{if } \text{dist}(\mathcal{C}) \geq 4, \end{cases}$$

density  $\Delta = V_{2n} \rho^{2n} 4^k 3^{-n/2}$ , where  $\rho = \frac{1}{2} d^{1/2}$ ,

$$\text{contact number } \tau = \begin{cases} 2^s A_s & \text{if } \text{dist}(\mathcal{C}) = \delta < 4, \\ 16A_4 + 6n & \text{if } \text{dist}(\mathcal{C}) = 4, \\ 6n & \text{if } \text{dist}(\mathcal{C}) > 4. \end{cases}$$

*Proof.* Suppose  $\text{dist}(\mathcal{C}) = 4$ . Then the lattice points closest to  $\mathbf{0}$  consist of  $16A_4$  {e.g. a codeword  $(1\omega 1\omega^2 0 \dots 0)$  gives rise to  $(1/2^{1/2})(\pm 1, \pm\omega, \pm 1, \pm\omega^2, 0, \dots, 0)$ } plus  $6n$  {of type  $((\pm\omega^r 2^{1/2})^1 0^{n-1})$  for  $r = 0, 1, 2$ }, for a total of  $\tau = 16A_4 + 6n$  at a squared distance of  $d = 2$ . Similarly for the other cases. The expression for  $\Delta$  follows from Eq. (1).

**THEOREM 7.**  $\Theta_{\Lambda(\mathcal{C})}(z) = W_{\mathcal{C}}(\phi_0(z), \phi_1(z))$ , where

$$\begin{aligned} \phi_0(z) &= \theta_2(2z) \theta_2(6z) + \theta_3(2z) \theta_3(6z) && \text{(see (13)),} \\ \phi_1(z) &= \theta_2(2z) \theta_3(6z) + \theta_2(6z) \theta_3(2z) \\ &= \frac{1}{2} \theta_2\left(\frac{1}{2}z\right) \theta_2\left(\frac{3}{2}z\right) \\ &= 2q^{1/2}(1 + q + 2q^3 + q^4 + 2q^6 + 2q^9 + \dots). \end{aligned} \tag{14}$$

*Proof.* Consider a codeword  $\mathbf{u} = (u_1, u_2, \dots) \in \mathcal{C}$ . The corresponding centers in  $\Lambda(\mathcal{C})$ , i.e. the points  $(1/2^{1/2}) \sigma^{-1}(\mathbf{u})$ , consist of the set

$$\begin{aligned} \Lambda(\mathbf{u}) &= \{(1/2^{1/2})(x_1, x_2, \dots): x_1 \in \sigma^{-1}(u_1), x_2 \in \sigma^{-1}(u_2), \dots\} \\ &= \{(y_1, y_2, \dots): y_r \in (1/2^{1/2}) u_r + 2^{1/2} \mathcal{E}, 1 \leq r \leq n\}. \end{aligned}$$

Let the theta-function of any subset  $\mathcal{S}$  of  $\Lambda(\mathcal{C})$  be

$$\Theta_{\mathcal{S}}(z) = \sum_{\mathbf{y} \in \mathcal{S}} q^{y \cdot \bar{y}}.$$

From (13),

$$\Theta_{2^{1/2}\mathcal{C}}(z) = \Theta_{\mathcal{C}}(2z) = \phi_0(z),$$

and a similar calculation shows that if  $\mathcal{S} = 2^{-1/2}w^r + 2^{1/2}\mathcal{E}$ ,  $r = 0, 1, 2$ , then

$$\Theta_{\mathcal{S}}(z) = \theta_2(2z) \theta_3(6z) + \theta_2(6z) \theta_3(2z) = \phi_1(z).$$

Hence

$$\begin{aligned} \Theta_{\Lambda(u)}(z) &= \phi_0(z)^{n-wt(u)} \phi_1(z)^{wt(u)}, \\ \Theta_{\Lambda(\mathcal{S})}(z) &= \sum_{u \in \mathcal{S}} \Theta_{\Lambda(u)}(z) = W_{\mathcal{S}}(\phi_0(z), \phi_1(z)). \end{aligned}$$

It remains to show that the first and second lines of (14) are equal. Using (12) this identity may be rewritten as

$$\theta_2(\tau) \theta_2(3\tau) + \theta_4(\tau) \theta_4(3\tau) = \theta_3(\tau) \theta_3(3\tau) \tag{15}$$

(apparently not in [41] or [43]). Then

$$\begin{aligned} &\theta_3(\tau) \theta_3(3\tau) - \theta_4(\tau) \theta_4(3\tau) \\ &= 2 \sum_{\substack{m=-\infty \\ m+n \text{ odd}}}^{\infty} \sum_{n=-\infty}^{\infty} q^{m^2+3n^2} \quad (\text{set } m = \frac{1}{2}(-a + 3b + 1), n = \frac{1}{2}(a + b + 1)) \\ &= 2 \sum_{\substack{a=-\infty \\ a+b \text{ odd}}}^{\infty} \sum_{b=-\infty}^{\infty} q^{(a+1/2)^2+3(b+1/2)^2} \\ &= \sum_{a=-\infty}^{\infty} \sum_{b=-\infty}^{\infty} q^{(a+1/2)^2+3(b+1/2)^2} = \theta_2(\tau) \theta_2(3\tau), \end{aligned}$$

since

$$\sum_{\substack{a=-\infty \\ a+b \text{ odd}}}^{\infty} \sum_{b=-\infty}^{\infty} q^{(a+1/2)^2+3(b+1/2)^2} = \sum_{\substack{c=-\infty \\ c+b \text{ even}}}^{\infty} \sum_{b=-\infty}^{\infty} q^{(-c-1/2)^2+3(b+1/2)^2}.$$

This completes the proof.

EXAMPLE 4. The trivial code  $T_n$  with generator matrix  $I_n$ ,  $k = n$ ,  $\text{dist} = 1$ ,  $W_{\mathcal{S}}(x, y) = (x + 3y)^n$  produces a lattice  $\Lambda(T_n)$  with  $\rho = 2^{-3/2}$ ,  $\Delta = V_{2n}(2 \cdot 3^{1/2})^{-n}$ ,  $\tau = 6n$ , and theta-function  $(\phi_0(z) + 3\phi_1(z))^n$ . This lattice is clearly equal to  $(1/2^{1/2}) \mathcal{E}^n$ , which has theta-function  $\phi_0(\frac{1}{2}z)^n$ , thus establishing the identity

$$\phi_0(z) + 3\phi_1(z) = \phi_0(\frac{1}{2}z). \tag{16}$$

(Equation (16) also follows from (12) and (15).) In particular when  $n = 1$  the corresponding real lattice  $\Lambda(T_1)_R$  is the densest packing  $A_2$  in  $\mathbb{R}^2$ .

EXAMPLE 2 (cont.). Similarly in  $\mathbb{C}^2$  the lattice  $\Lambda(C_2)$  is equal to  $\mathcal{E}^2$ , which has theta-function  $\phi_0(\frac{1}{2}z)^2$ , leading to the identity

$$\phi_0(z)^2 + 3\phi_1(z)^2 = \phi_0(\frac{1}{2}z)^2.$$

EXAMPLE 5. In  $\mathbb{C}^3$  three copies of  $\Lambda(T_3)$  can be fitted together without overlap. The resulting lattice is

$$\Lambda(T_3) \cup (\mathbf{u} + \Lambda(T_3)) \cup (2\mathbf{u} + \Lambda(T_3))$$

where  $\mathbf{u} = 2^{1/2} 3^{-1}(\lambda, \lambda, \lambda)$ ,  $\lambda = 1 - \omega$ ; and has generator matrix

$$\frac{1}{2^{1/2} 3} \begin{bmatrix} 3 & 0 & 0 \\ 0 & 3 & 0 \\ \lambda & \lambda & \lambda \end{bmatrix},$$

$\rho = 2^{-3/2}$ ,  $\Delta = V_6/8(3^{1/2})$ , and  $\tau = 72$ . The corresponding real lattice is  $E_6$  (cf. [5]) and in fact this is equivalent to Coxeter's construction of  $E_6$  (see for example [9, p. 421]). It can be seen that the theta-function of  $E_6$  is

$$\phi_0(\frac{1}{2}z)^3 + \frac{1}{4}\{\phi_0(\frac{1}{2}z) - \phi_0(\frac{1}{4}z)\}^2.$$

EXAMPLE 3 (cont.). In  $\mathbb{C}^6$  we find that  $\Lambda(Q_6)$  is a self-dual lattice ( $U_6$  in Feit's notation [14a]) with  $\Delta = V_{12}/27$  and  $\tau = 756$ . The corresponding real lattice is  $K_{12}$ , the densest known packing in  $\mathbb{R}^{12}$ , and in fact this is equivalent to Coxeter and Todd's construction of  $K_{12}$  given in [11]. The automorphism group of  $\Lambda(Q_6)$  is  $[2 \ 1; 3]^3$ , a unitary group generated by reflections (see [16, 36, 37, 42]).

The theta-function  $\Theta_{\Lambda(Q_6)}(z)$  is given by Theorem 7, but we are more interested in the cusp form  $\Delta_6(z)$ , defined by

$$\begin{aligned} \Delta_6(z) &= \frac{1}{36} \{ \Theta_{\Lambda(C_2)}(z)^3 - \Theta_{\Lambda(O_6)}(z) \} \\ &= \frac{1}{36} \{ (\phi_0(z)^2 + 3\phi_1(z)^2)^3 - (\phi_0(z)^6 + 45\phi_0(z)^2 \phi_1(z)^4 + 18\phi_1(z)^6) \} \\ &= \frac{1}{4} \phi_1(z)^2 \{ \phi_0(z)^2 - \phi_1(z)^2 \}^2. \end{aligned}$$

From (10), (11), (15),

$$\begin{aligned} \phi_1(x)^2 &= \frac{1}{2} \{ \theta_2(x)^2 \theta_2(3x)^2 + \theta_3(x)^2 \theta_3(3x)^2 - \theta_4(x)^2 \theta_4(3x)^2 \} \\ &= \theta_2(x) \theta_2(3x) \theta_3(x) \theta_3(3x), \end{aligned} \tag{17}$$

$$\begin{aligned} \phi_0(x)^2 &= \frac{1}{2} \{ \theta_2(x)^2 \theta_2(3x)^2 + \theta_3(x)^2 \theta_3(3x)^2 + \theta_4(x)^2 \theta_4(3x)^2 \}, \\ \phi_0(x)^2 - \phi_1(x)^2 &= \theta_4(x)^2 \theta_4(3x)^2. \end{aligned} \tag{18}$$

Therefore from (10), (17), (18), (8) and (7),

$$\begin{aligned} \Delta_6(z) &= \frac{1}{16} \left\{ \theta_2 \left( \frac{z}{2} \right) \theta_2 \left( \frac{3z}{2} \right) \theta_3 \left( \frac{z}{2} \right) \theta_3 \left( \frac{3z}{2} \right) \theta_4 \left( \frac{z}{2} \right) \theta_4 \left( \frac{3z}{2} \right) \right\}^2 \\ &= \frac{1}{16} \theta'_1 \left( \frac{z}{2} \right)^2 \theta'_1 \left( \frac{3z}{2} \right)^2 \\ &= q \prod_{m=1}^{\infty} (1 - q^m)^6 (1 - q^{3m})^6. \end{aligned} \tag{19}$$

The coefficients of  $\Delta_6(z)$  play the role of the Ramanujan numbers (see Theorem 9). The first few are as follows.

$$\begin{aligned} \Delta_6(z) &= q - 6q^2 + 9q^3 + 4q^4 + 6q^5 - 54q^6 - 40q^7 \\ &\quad + 168q^8 + 81q^9 - 36q^{10} - 564q^{11} + 36q^{12} + 638q^{13} \\ &\quad + 240q^{14} + 54q^{15} - 1136q^{16} + 882q^{17} - 486q^{18} \\ &\quad - 556q^{19} + 24q^{20} + \dots \end{aligned} \tag{20}$$

### 7. A BASIS FOR THE THETA-FUNCTIONS

Let  $\mathcal{M}_n$  be the complex vector space spanned by the theta-functions of all self-dual  $\mathcal{E}$ -lattices in  $\mathbb{C}^n$ , and let  $\mathcal{M} =: \bigoplus_{n=0}^{\infty} \mathcal{M}_n$ , a graded ring.

**THEOREM 8.**

$$\dim_{\mathbb{C}} \mathcal{M}_n = 1 + \left\lfloor \frac{n}{6} \right\rfloor, \tag{21}$$

$$\sum_{n=0}^{\infty} \lambda^n \dim_{\mathbb{C}} \mathcal{M}_n = \frac{1}{(1 - \lambda)(1 - \lambda^6)}. \tag{22}$$

*Proof.* If  $\Theta(z) \in \mathcal{M}_n$  let  $\Psi(z) = \Theta(2z/3^{1/2})$  so that from Corollary 2

$$\Psi(z + 3^{1/2}) = \Psi(z), \quad \Psi\left(-\frac{1}{z}\right) = \left(\frac{z}{i}\right)^n \Psi(z).$$

Hence  $\Psi(z) \in \mathcal{M}(3^{1/2}, n, 1)$ , in the notation of Ogg [32, p. xiv]. Then (21) follows from Ogg [32, p. I-23, Theorem 3 with  $q = 6$ ], and (22) is immediate.

**THEOREM 9.**  $\mathcal{M} = \mathbb{C}[\Theta_{\mathcal{E}}(z), \Delta_6(z)]$  where  $\Theta_{\mathcal{E}}(z), \Delta_6(z)$  are given by (13), (19). Therefore the theta-function of a self-dual  $\mathcal{E}$ -lattice in  $\mathbb{C}^n$  can be written

$$\Theta_{\Lambda}(z) = \sum_{r=0}^{\lfloor n/6 \rfloor} c_r \Theta_{\mathcal{E}}(z)^{n-6r} \Delta_6(z)^r, \tag{23}$$

for uniquely determined integers  $c_r$ .

*Proof.* Since  $\Theta_g(z) \in \mathcal{M}_1$  and  $\Delta_6(z) \in \mathcal{M}_6$  are algebraically independent this follows immediately from Eq. (22).

*Remarks.* (i) Equation (6) may be verified for  $\Theta_g(z)$  and  $\Delta_6(z)$  using (8), (9) and (12).

(ii) A modification of [32, p. I-32] shows that  $E_6(3z/2) - \frac{1}{2}E_6(z/2) \in \mathcal{M}_6$ , where  $E_6(z)$  denotes the Eisenstein series  $1 - 504\sum\sigma_5(m)q^{2m}$  ([32, p. I-28; 35, p. 93]). This leads to the identity

$$27E_6\left(\frac{3z}{2}\right) - E_6\left(\frac{z}{2}\right) = 26\Theta_g(z)^6 - 432\Delta_6(z). \tag{24}$$

Theorem 9 is an analog of the classical result that the theta-function of an even unimodular lattice is an element of  $\mathbb{C}[E_4(z), \Delta(z)]$ , where  $E_4(z)$  is an Eisenstein series and  $\Delta(z) = q^2 \prod_{m=1}^{\infty} (1 - q^{2m})^{24} = \sum_{m=0}^{\infty} a_m q^{2m}$ , the  $a_m$  being the Ramanujan numbers (see for example [6, Theorems 5, 6; 7; or 40]). Just as the Ramanujan numbers are multiplicative [15, p. 68, Corollary 5; 32, p. II-15] so are the coefficients of  $\Delta_6(z)$ . More precisely, let  $\Delta_6(z) = \sum_{m=1}^{\infty} a_m q^m$  and let  $\phi_6(s)$  be the Dirichlet series

$$\phi_6(s) = \sum_{n=1}^{\infty} \frac{a_n}{n^s}, \quad \text{Re}(s) > 4. \tag{25}$$

**THEOREM 10.**  $\phi_6(s)$  has an Euler product

$$\phi_6(s) = (1 - 9 \cdot 3^s)^{-1} \prod_{p \neq 3} (1 - a^p p^{-s} + p^{5-2s})^{-1}, \tag{26}$$

and therefore the multiplication rule for the  $a_n$  is (compare (20))

$$a_m a_n = \prod_{3 \nmid d, d|(m,n)} d^5 a_{mn/d^2}. \tag{27}$$

*Proof.* From Corollary 2,  $\Delta_6(2z)$  is the unique cusp form of weight 6 for  $\Gamma_0(3)$ . From [3, Lemma 27; 38, Sect. 3.5],  $\Delta_6(2z)$  is an eigenfunction for the Hecke operators  $T'(n)_{6,\psi}$  for all  $n$ , where  $\psi(m) = 0$  if  $3 \mid m$ ,  $\psi(m) = 1$  if  $3 \nmid m$ . Then [38, Theorem 3.43] implies (26).

Suppose the coefficients  $c_0, \dots, c_\mu$ ,  $\mu = [n/6]$ , in (23) are chosen so as to make the  $q$ -expansion of the right-hand side begin

$$1 + A^*_{\mu+1} q^{\mu+1} + A^*_{\mu+2} q^{\mu+2} + \dots. \tag{28}$$

The resulting series is called the *extremal theta-function* in dimension  $n$  (cf. [30]), and a lattice with this theta-function (if one exists) is called an *extremal lattice*. The examples  $\mathcal{E}^n$  ( $1 \leq n \leq 5$ ) and  $\Lambda(Q_6)$  are extremal lattices in dimensions 1-6. Examples for  $n = 8, 10$  are provided by the lattices  $\Lambda(E_8)$  (or Example 6

below) and  $\Lambda(E_{10})$  (or  $\Lambda(B_{10})$ ), using codes defined in [27]. For other examples with  $n \leq 12$  see [14a]. The methods of [30] may be applied here, and show among other things that:

**THEOREM 11.** *For all  $n$ ,  $A^*_{\mu+1} > 0$ , and so a self-dual  $\mathcal{E}$ -lattice has  $d \leq 1 + [n/6]$ . However, for all sufficiently large  $n$ ,  $A^*_{\mu+2} < 0$ , and no extremal self-dual lattice exists in  $\mathbb{C}^n$ .*

*Remark.* The complex version of the Leech lattice is an  $\mathcal{E}$ -lattice in  $\mathbb{C}^{12}$  with  $\Delta = V_{24}$ ,  $\tau = 196560$ , but is not self-dual [8, 14a, 23].

### 8. CONSTRUCTION B

Let  $\mathcal{C}$  be a linear code over  $GF(4)$  of length  $n$ , dimension  $k$ , and minimum distance 8, satisfying  $\mathcal{C} \subset \mathcal{C}^\perp$  and such that the all-ones vector is in  $\mathcal{C}$ . The latter is not a serious restriction in view of Corollary 9 of [27]. These assumptions imply that if  $\mathbf{u} \in \mathcal{C}$  contains  $i$  1's,  $j$   $\omega$ 's and  $k$   $\omega^2$ 's, then  $i \equiv j \equiv k \equiv 0 \pmod{2}$ .

#### Construction B

With  $\mathcal{C}$  as above, let  $\mathcal{L}(\mathcal{C})$  consist of all  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{C}^n$  such that (i)  $\sigma(\mathbf{x}) \in \mathcal{C}$ , and (ii)  $\sum_{i=1}^n x_i \in 4\mathcal{E}$ .

The following facts are now easily verified:  $\mathcal{L}(\mathcal{C})$  is an  $\mathcal{E}$ -lattice in  $\mathbb{C}^n$  with  $\det \mathcal{L}(\mathcal{C}) = 2^{n-k+1}$ ,  $\rho = 2^{1/2}$ ,  $\Delta = V_{2n} 2^{2k-23-n/2}$ , and  $\tau = 192A_8^{(b)} + 64A_8 + 6n(n-1)$ , where  $A_8$  is the total number of codewords in  $\mathcal{C}$  of weight 8, and  $A_8^{(b)}$  is the number with 8 1's and  $(n-8)$  0's. There is a formula for the theta-function of  $\mathcal{L}(\mathcal{C})$ , analogous to but more complicated than Theorem 7. We omit the details.

**EXAMPLE 6.** The code  $\{0^8, 1^8, \omega^8, (\omega^2)^8\}$  produces an extremal lattice in  $\mathbb{C}^8$  with  $\tau = 720$  and  $\Delta_{16} \cdot 3^{-4}$ .

**EXAMPLE 7.** The code  $S_{18}$  given in [27] produces a lattice in  $\mathbb{C}^{18}$  with  $\Delta = V_{36} 2^{163-9} \approx 3.330V_{36}$ . The corresponding real lattice in  $\mathbb{R}^{36}$  also has  $\Delta \approx 3.330V_{36}$ , and appears to be the densest lattice yet found in this dimension. On the other hand  $\tau < 189000$ , compared with the present record of 234456 [21]. Incidentally an extremal lattice in  $\mathbb{C}^{18}$  would have  $\tau = 646380$ .

*Remark.* One may also define the coordinate array of a point in  $\mathcal{E}^n$ , and hence generalize Construction C of [20, 21, 39] to  $\mathcal{E}$ -lattices.

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